Reusable formal models for secure software architectures

Thomas Heyman, Riccardo Scandariato, Wouter Joosen

IBBT-DistriNet, KU Leuven

Problem

Formal methods for secure software architectures

Enforce rigor Enable reasoning Provide assurance → interesting for SA

Formal methods not widely used...

High overhead Require **expertise** Different **stakeholders**

Solution

Contribution – part I for the security engineer



Refined models

of building blocks (e.g., security patterns)

Created and used by security engineer → Assess security Results in better documentation Verification results are reusable

Contribution – part II for the software architect



Abstract models

of building blocks (e.g., security patterns)

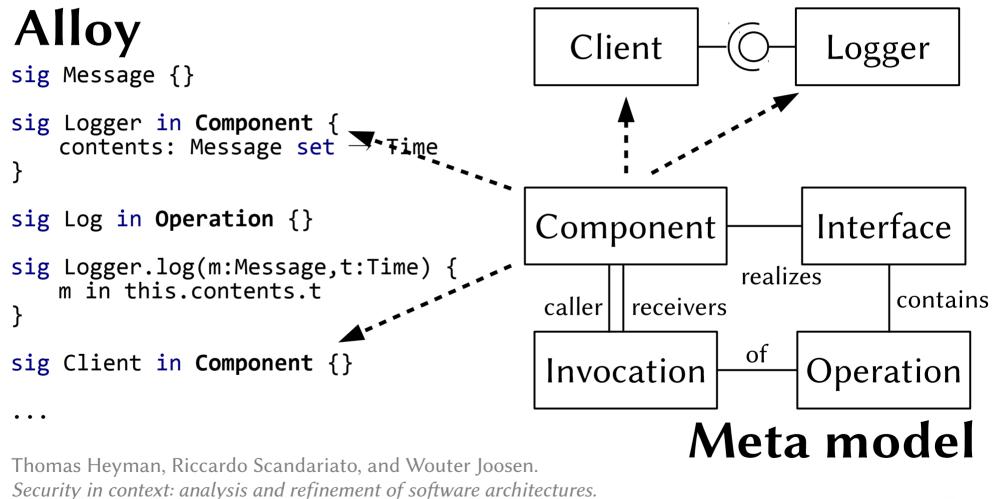
Simple, behaves like refinement **Created** by security engineer **Used** by software architect

→ Uncover compositional issues (Re-)usable!

Outline Background Contribution I Contribution II Wrap-up

Background modelling software architectures

Architecture



In Annual IEEE Computer Software and Applications Conference, July 2010.



Modelling a pattern language for accountability

A pattern language for accountability

Contents

Secure Logger, Audit Interceptor, Authentic. and Authoriz. Enforcer,

Secure Pipe Christopher Steel, Ramesh Nagappan, and Ray Lai.

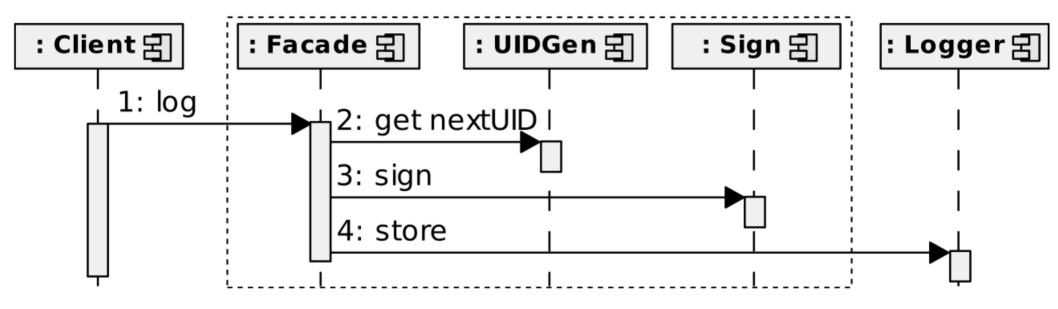
Christopher Steel, Ramesh Nagappan, and Ray Lai. Core Security Patterns: Best Practices and Strategies for J2EE, Web Services, and Identity Management. Prentice Hall, 2005.

Motivation

Self-contained set

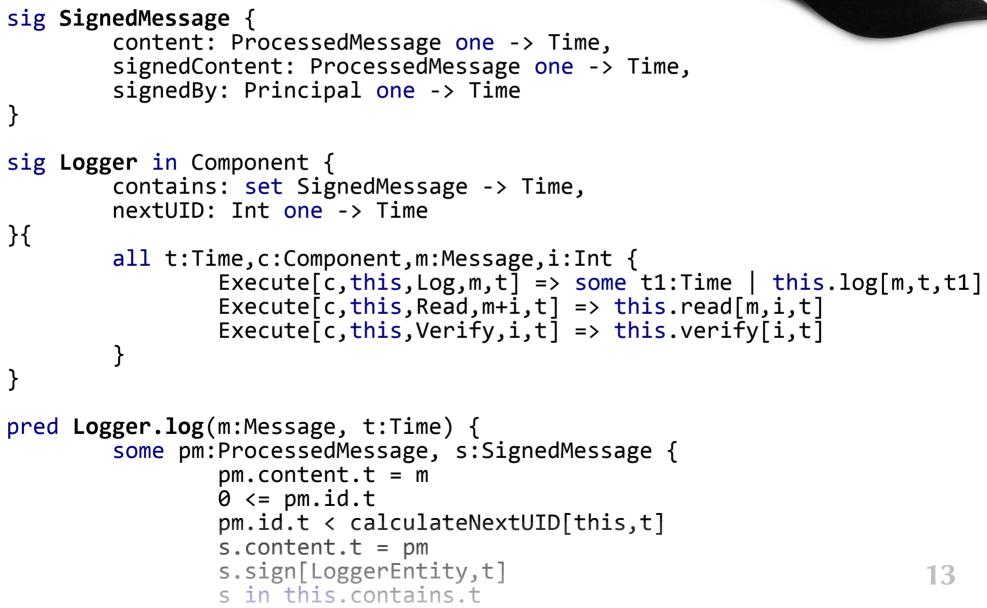
Useful in practice (industrial projects)

Modelling the Secure Logger pattern



Christopher Steel, Ramesh Nagappan, and Ray Lai. *Core Security Patterns: Best Practices and Strategies for J2EE, Web Services, and Identity Management.* Prentice Hall, 2005.

Modelling the Secure Logger pattern



Verification encoding sec. requirements

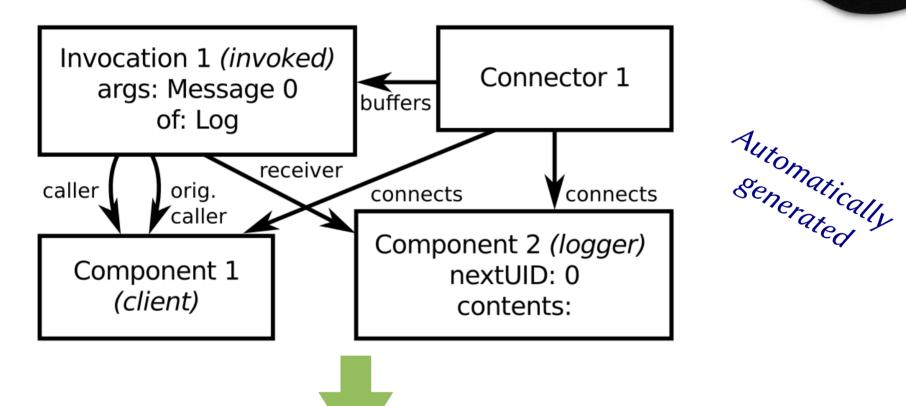


"Whenever a message is logged, it can be read back later or the verify method returns *false*."

assert NothingDeleted {
all t:Time,m:Message,l:Logger,c:Component |
 Invoke[c,l,Log,m,t] implies (
 some t1:t.nexts+t {

Verification analyzing counterexamples





"assume that invocations are eventually executed"



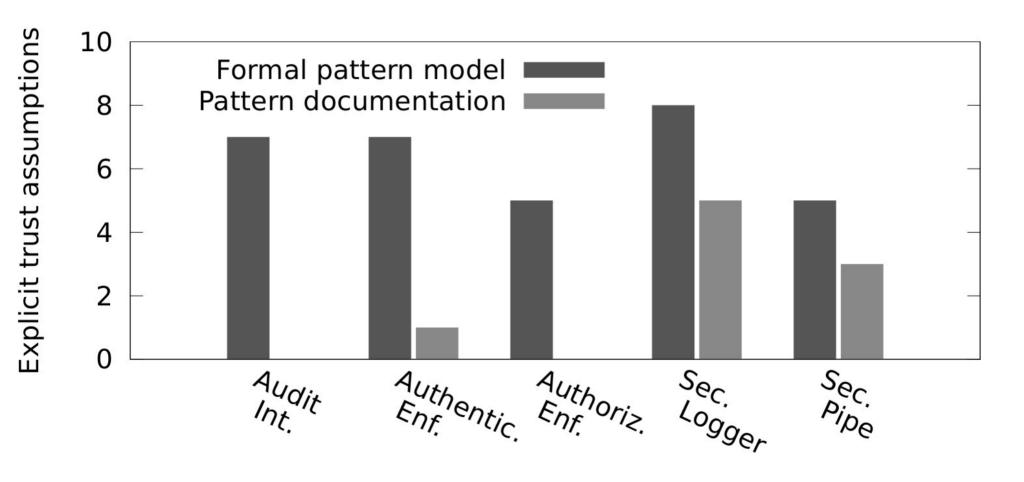
Contribution 1



Trust assumptions! Usually left implicit Assurance requires explicit assumptions

Modelling and verification... Makes them explicit Finds extra assumptions

Uncovered assumptions





Composing abstract models

Contribution 2

Abstraction

pred Logger.log(m:Message,t:Time) {
 some c:Component,t1:t.prevs+t | Execute[c,this,Log,m,t1]



VS.



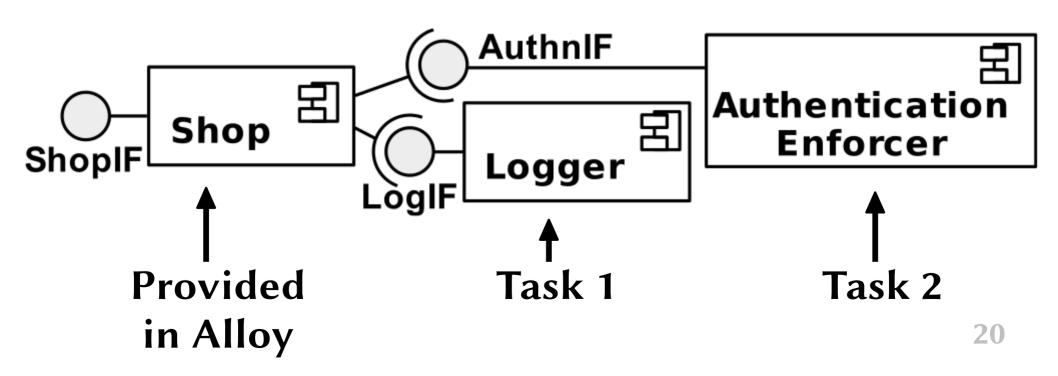
Refinement

pred Logger.log(m:Message, t:Time) {
 some pm:ProcessedMessage, s:SignedMessage {
 pm.content.t = m
 0 <= pm.id.t and pm.id.t < calculateNextUID[this,t]
 s.content.t = pm and s.sign[LoggerEntity,t]
 s in this.contains.t</pre>

Case study

Two subjects senior researchers

Extend basic architecture



Case study results

Both candidates successful 1St 1 hour, 2nd 2 ½ hour + exit questionnaire = useful

Results both solutions correct (in line with reference solution) ±7 assumptions each 1 flaw in solution, results in assumption

Summary what to take home...

Modelled pattern language for accountability

Verify Once, Reuse Many

Provides insight in patterns

Big picture this research in context

Larger research track Formal methods in secure software architecture

Under review: formal framework In progress: DSL + tool support

http://distrinet.cs.kuleuven.be/software/samodels/ thomas.heyman@cs.kuleuven.be